

1. Solve the recurrence $U(t) = ab^t + cU(t - 1)$.
2. Approximate $\cos x - 1 - x^2/2$ with accuracy of one term plus a big O .
3. You have n tapes, numbered from 0 to $n - 1$. Each tape contains one sorted block. Each block has the same number of items. You also have a supply of empty tapes. You wish to produce a tape containing one sorted block by merging tapes 0 to $n - 1$. Assume that one unit of time is needed for each item output.
 - a. One way you can do this is merging tapes i and $n - 1 + i$ to produce tape $n + i$, starting with $i = 0$ and continuing until all the items have been merged. After the first round, this algorithm merges the tapes one at a time into the previous result. How many time units does this method take?
 - b. Another way to do this is by merging tapes $2i$ and $2i + 1$ to produce tape $n + i$. This algorithm merges equal size tapes in pairs until the final result is obtained. How many time units does this method take?
 - c. Which method is faster? Is the difference significant? Give an informal explanation of the results of this question.

4. Let

$$p(k, i) = \binom{v-i}{k} (2p)^k (1-p)^{2v-i-k},$$

and

$$s(k, m, j, i) = \frac{\binom{k}{j} \binom{v-i-k}{m-j}}{\binom{v-i}{m} 2^j}.$$

Simplify

$$\sum_j p(j+k, i) s(j+k, m, j, i).$$

5. Suppose

$$v^n \leq 1 + 2 \left[1 + O\left(\frac{1}{pv}\right) \right] [1 + 2pv(1 + p^2)v]^t,$$

where all variables are positive, n is fixed, v goes to infinity, t goes to infinity, p goes to zero, and pv goes to infinity. Convert this inequality into a condition on p . (In other words solve for p .)