

PhD Qualifying Exam in Foundations

August 23, 2005

1. Show directly that if a language is defined by a regular expression then it is context-free.
2. Let M be a deterministic finite automaton over the alphabet $\{a, b\}$. Show that the following language is regular:

$\{w \in \{a, b\}^* \mid \text{the run of } M \text{ on } w \text{ goes through each accepting state of } M \text{ at least once}\}$

3. Consider the language L over the alphabet $\{a, b\}$ defined by

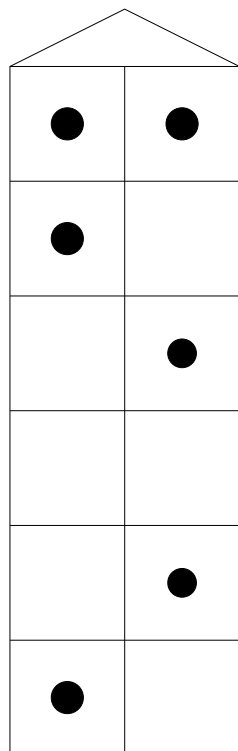
$$L = \{a^n b a^{n+1} \mid n \geq 0\}.$$

- (a) Give a context-free grammar that generates L . Prove that every string in L is generated by your grammar (you need not prove the converse implication).
 - (b) Prove that L is not regular. For full credit use the Myhill-Nerode Theorem. For half credit use the Rabin-Scott Pumping Lemma.
4. For each of the following problems, state whether it is decidable and either provide a decision algorithm, or prove that none exists. You may quote major theorems without proof.
 - (a) Does a given deterministic finite automaton accept all strings over its alphabet?
 - (b) Is the language accepted by a given Turing machine regular?

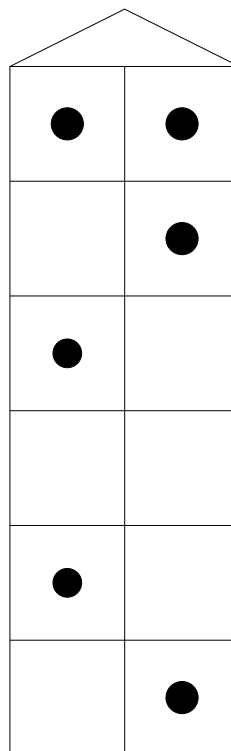
5. Let P be the collection of languages decidable in polynomial time. Under which of the following operations is P closed: (a) intersection; (b) complementation; (c) Kleene's star. Prove your answers succinctly.

6. Give an effective enumeration M_1, M_2, \dots of some Turing machines where the collection of these machines has the following two properties:
 - (a) Each language $L(M_i)$ is NP.
 - (b) Each NP language is $L(M_i)$ for some i .

7. Consider the following collection of puzzles. For $m \geq 1$ and $n \geq 1$, an m - n puzzle consists of a deck of m cards, each displaying a $2 \times n$ grid, with the “top” of the card identified by a triangular cap. Each square of the grid may have a circular hole punched at its center. Here is an example of a possible card in a m - n puzzle where $n = 6$:



A 6-puzzle card



Same card flipped

A deck of such cards may have a “see-through” hole, when a certain position of the grid has a hole on all the cards. A puzzle is “solvable” if it is possible to flip some of the cards horizontally (flip 180 degrees along the vertical axis, as shown in the figure) to yield a deck with no see-through holes in the *left* column.

Show that the collection of solvable m - n puzzles is NP-complete.